

Rare event probability estimation for connectivity of large random graphs

Rohan Shah

School of Mathematics and Physics
The University of Queensland
Brisbane QLD 4702, AUSTRALIA

Christian Hirsch

Institute of Stochastics
Ulm University
89069 Ulm, GERMANY

Dirk P. Kroese

School of Mathematics and Physics
The University of Queensland
Brisbane QLD 4702, AUSTRALIA

Volker Schmidt

Institute of Stochastics
Ulm University
89069 Ulm, GERMANY

ABSTRACT

Spatial statistical models are of considerable practical and theoretical interest. However, there has been little work on rare-event probability estimation for such models. In this paper we present a conditional Monte Carlo algorithm for the estimation of the probability that random graphs related to Bernoulli and continuum percolation are connected. Numerical results are presented showing that the conditional Monte Carlo estimators significantly outperform the crude simulation estimators.

1 INTRODUCTION

Random graph models are of significant practical importance; see, e.g., Sahini and Sahimi 1994. The connectivity properties of such models are of considerable interest, for example in network reliability (Gertsbakh and Shpungin 2010, Colbourn 1987), percolation theory (Bollobás and Riordan 2006) and material design (Stenzel, Koster, Thiedmann, Oosterhout, Janssen, and Schmidt 2012). In percolation theory the focus is on infinite random graph models, which are theoretically more tractable. However physical systems of interest are necessarily finite and this suggests the use of finite random graph models in applications.

This paper studies the following problem: consider a connected ‘base’ graph \mathcal{G} , and retain vertices independently with probability p . If we use this random vertex subset to construct the induced subgraph, what is the probability that the induced subgraph is connected? Calculating this probability exactly for a finite but large random graph model constitutes a difficult counting problem. In the network reliability setting this problem has been proved to be #P-complete (Colbourn 1987). Given the difficulties with exact computation we naturally turn to Monte Carlo methods.

However, crude Monte Carlo techniques can be cumbersome because connectivity is often a rare event and the problem becomes one of rare-event simulation. This is similar to the situation in network reliability, which also involves rare-event simulation; however in that case disconnection is the rare event, rather than connection. Typical methods for efficient rare event simulation include splitting (Kahn and Harris 1951, Glasserman, Heidelberger, Shahabuddin, and Zajic 1999, Garvels, van Ommeren, and Kroese 2002, L’Ecuyer, Demers, and Tuffin 2006, Botev and Kroese 2012), importance sampling (Glynn and Iglehart 1989, Asmussen and Rubinstein 1995) and conditional Monte Carlo (Asmussen and Glynn 2007). See Rubinstein and Kroese 2008 or Kroese, Taimre, and Botev 2011 for an overview of these techniques.

Random *geometric* graphs are the continuous analog of random graph models. The defining property of these models is that the vertices of the graph are the points of a point process on a bounded sampling window. Although these models can be viewed as strictly combinatoric, the spatial structure of the model remains important. Even in the infinite domain case very little has been proved about the connectivity properties of such models and related critical exponents (Brereton, Hirsch, Kroese, and Schmidt 2014). This has led to the widespread use of Monte Carlo methods to estimate unknown percolation thresholds (Quintanilla and Ziff 2007, Li and Östling 2013, Torquato and Jiao 2012). We show that the Monte Carlo estimate we propose can be applied to the Gilbert disk model with minimal change.

The rest of this paper is organized as follows. Section 2 describes the Bernoulli site percolation model on a finite lattice and the Gilbert disk model. Section 3 outlines the conditional Monte Carlo estimator for the Bernoulli site percolation model. Section 4 describes the adaptation of the conditional Monte Carlo estimator in Section 3 to the Gilbert disk model. Section 5 gives numerical results showing that the conditional Monte Carlo estimators perform significantly better than the crude simulation estimators.

2 PRELIMINARIES

Let $\mathcal{G} = (\mathcal{V}, \mathcal{E})$ be a finite connected graph. For any vertex v the *degree* of v is the number of edges incident to v , and we write $\deg(v)$. The maximum degree of any vertex in \mathcal{G} is denoted by $\Delta(\mathcal{G})$. The cardinality of a finite set S is denoted by $|S|$. If S is uncountable then $|S|$ denotes instead the Lebesgue measure of the set.

Take some $p \in (0, 1)$ and let $q = 1 - p$. Let $X = \{X_v\}_{v \in \mathcal{V}}$ be a collection of independent and identically distributed (iid) random variables with $X_v \sim \text{Ber}(p)$. A vertex v with $X_v = 1$ is called *activated*. Let $V(X)$ denote the set of activated vertices, that is

$$V(X) = \{v \in \mathcal{V} \mid X_v = 1\}.$$

The random subset $V(X)$ induces a random subgraph

$$G = G(X) = (V(X), E(X)),$$

where

$$E = E(X) = \{(v_1, v_2) \in \mathcal{E} \mid v_1, v_2 \in V(X)\}.$$

We typically omit the dependence of these random variables on X . We denote the collection of possible induced subgraphs of \mathcal{G} by $\mathcal{P}(\mathcal{G})$. We will write the density of G with respect to counting measure on $\mathcal{P}(\mathcal{G})$ as $f_G(g; p)$. For $g_1, g_2 \in \mathcal{P}(\mathcal{G})$ induced by vertex subsets $V_1, V_2 \subseteq \mathcal{V}$ we will write $g_1 \cap g_2$ for the subgraph induced by the vertex set $V_1 \cap V_2$.

Models of this form for G are commonly known as *discrete site percolation models*, although \mathcal{G} is often implicitly assumed to be infinite. We will also refer to the case where \mathcal{G} is an arbitrary finite graph as being a discrete site percolation model. As vertices are retained independently with some probability p these models are said to be *Bernoulli site percolation models*.

Let $\mathcal{C} \subseteq \mathcal{V}$ be a subset of vertices such that the subgraph induced by the subset is *connected*. Define $\partial\mathcal{C}$ to be the *boundary vertices* of \mathcal{C} in \mathcal{G} . That is,

$$\partial\mathcal{C} = \{v_1 \in \mathcal{V} \mid v_1 \notin \mathcal{C}, \{v_1, v_2\} \in \mathcal{E} \text{ for some } v_2 \in \mathcal{C}\}.$$

Define the connectivity probability

$$\ell(\mathcal{G}, p) = \mathbb{P}(G \text{ is connected}).$$

If we add vertices to G while maintaining a bound on the maximum vertex degree $\Delta(\mathcal{G})$, then $\ell(\mathcal{G}, p)$ will decay exponentially fast in the number of vertices. See Weichenberg, Chan, and Medard 2004 for

results bounding the connectivity probability in the related network reliability setting. The idea of *prime failure events* used in Weichenberg, Chan, and Medard 2004 applies equally to our site-percolation model. Exponential decay means in particular that $\ell(\mathcal{G}, p)$ will be small for large base graphs \mathcal{G} when $\Delta(\mathcal{G})$ is small.

The defining property of *random geometric graphs* is that their vertices are the points of a spatial point process on some bounded Borel set $R \subseteq \mathbb{R}^d$. Edges are added between vertices according to some probabilistic or deterministic rule. One possibility is to connect each vertex to the k closest other vertices; another is to connect a pair of vertices with some probability that depends on the Euclidean distance between them.

We focus on the *standard Gilbert disk model*, a special case of the *Boolean model* (Chiu, Stoyan, Kendall, and Mecke 2013). In this model the point process ξ that generates the vertices of the graph is a homogeneous Poisson point process on R with some intensity $\lambda > 0$, and any pair of vertices that are closer than some fixed distance r are connected by an edge. We will denote this model by $G_{\text{geo}}(R, \lambda, r)$, generally abbreviated to G_{geo} . The open ball of radius r around a point $x \in \mathbb{R}^d$ will be denoted by $B(x, r)$.

3 CONDITIONAL MONTE CARLO FOR DISCRETE PERCOLATION

If $\{X^{(i)}\}_{i=1}^{\infty}$ are iid copies of X then the *crude simulation estimator* is

$$\widehat{\ell}_{\text{crude}}(\mathcal{G}, p) = \frac{1}{n} \sum_{i=1}^n \mathbb{I} \left\{ G(X^{(i)}) \text{ is connected} \right\}, \quad (1)$$

where $n \geq 1$ is an arbitrary fixed integer and $\mathbb{I}\{A\}$ denotes the indicator function of an event A . Our aim is to find an estimator that has better asymptotic properties than the crude simulation estimator, as the number of vertices in \mathcal{G} is allowed to increase.

We can construct a simple conditional Monte Carlo estimator based on knowledge of a single connected component. After this connected component has been generated it is no longer necessary to simulate the states of the remaining vertices, as the connectivity probability can be computed exactly; it is the probability that the vertices not already simulated are all deactivated. By the total variance formula this gives an estimate with smaller variance than the crude estimator given in Equation (1). See Billingsley 1995, p. 275,445 for further details.

The idea of the algorithm is as follows. Select vertices randomly without replacement and generate their activation state according to a $\text{Ber}(p)$ -distributed random variable. Continue this process until some activated vertex is selected. The set of deactivated vertices is denoted by V_{deact} , with both $V_{\text{deact}} = \emptyset$ and $V_{\text{deact}} = \mathcal{V}$ being possible.

If an activated vertex is generated denote it by ω . We can then simulate the entire connected component C_{act} for ω by performing a depth-first search of $\mathcal{V} \setminus V_{\text{deact}}$. For every visited vertex v the activation state is generated according to a $\text{Ber}(p)$ -distributed random variable, and if v is activated the search continues to the neighbors of v . If no activated vertex was originally found, set $C_{\text{act}} = \emptyset$. The random object we will condition on is $Z_{\text{Ber}} = (V_{\text{deact}}, C_{\text{act}})$. It will be convenient to define $N_{\text{deact}} = |V_{\text{deact}}|$. Note that the random variables defined in this section are not just functions of the binary vector X .

The process of generating Z_{Ber} is illustrated in Figure 1. In this case V_{deact} contains the first three vertices selected, all of which were generated to be deactivated. The fourth vertex picked is marked as ω , and was determined to be activated. The connected component for ω was then generated, and contains 3 vertices. Note that v_1 was already determined to be deactivated when we started to generate C_{act} . The activation state has only been generated for the marked vertices; the activation states of the unmarked vertices is unknown.

The density of $G|Z_{\text{Ber}}$ is

$$f_{G|Z_{\text{Ber}}}(g|(v_{\text{deact}}, c_{\text{act}}); p) = \binom{|\mathcal{V}| - |k|}{|g| - |c_{\text{act}}|} p^{|g| - |c_{\text{act}}|} q^{(|\mathcal{V}| - |k|) - (|g| - |c_{\text{act}}|)}, \quad (2)$$

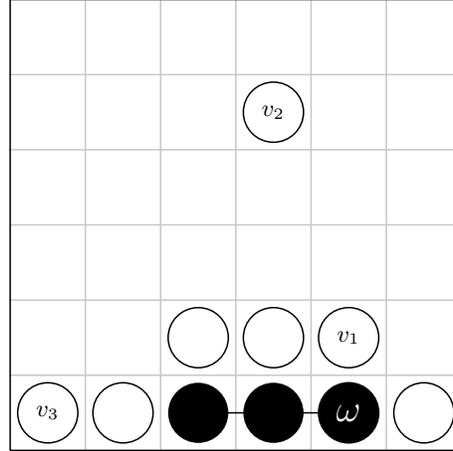


Figure 1: An example showing the process of generating Z_{Ber} where \mathcal{G} is the 6×6 grid graph which does not include diagonal edges. C_{act} contains three vertices. V_{deact} also contains three vertices, labeled v_1 , v_2 and v_3 . Activated vertices are marked by filled circles and deactivated vertices by empty circles. If a vertex is not shown then its state is still unknown.

where $k = v_{\text{deact}} \cup \partial c_{\text{act}} \cup c_{\text{act}}$. If $f_N(x; n, p)$ denotes the density of a Binomial (n, p) -distributed random variable then we can rewrite (2) as

$$f_{G|Z_{\text{Ber}}}(g | (v_{\text{deact}}, c_{\text{act}}); p) = f_N(|g| - |c_{\text{act}}|; |\mathcal{V}| - |k|, p).$$

Conditional on Z_{Ber} , the only way for G to be connected is if all vertices outside the set $K = V_{\text{deact}} \cup \partial C_{\text{act}} \cup C_{\text{act}}$ are deactivated. This has probability

$$\begin{aligned} \mathbb{P}(G \text{ is connected} \mid Z_{\text{Ber}} = z_{\text{Ber}}) &= \mathbb{P}(\mathcal{V} \setminus (c_{\text{act}} \cup \partial c_{\text{act}} \cup v_{\text{deact}}) \text{ are deactivated}) \\ &= q^{|\mathcal{V}| - |c_{\text{act}} \cup \partial c_{\text{act}} \cup v_{\text{deact}}|}. \end{aligned} \quad (3)$$

Note that if $V_{\text{deact}} = \mathcal{V}$ then G is the empty graph which is considered connected. The conditional probability is simple to calculate, and Z_{Ber} is simple to simulate. We can use the expression

$$\ell(\mathcal{G}, p) = \mathbb{E}[\mathbb{P}(G \text{ is connected} \mid Z_{\text{Ber}})] \quad (4)$$

to construct the following conditional Monte Carlo estimator for the Bernoulli site percolation model.

Proposition 1 (Conditional Monte Carlo estimator for the Bernoulli site percolation model)

Let $\{Z_{\text{Ber}}^{(i)}\}_{i=1}^{\infty}$ be iid copies of Z_{Ber} and $\{X^{(i)}\}_{i=1}^{\infty}$ be iid copies of X , where each of the $Z_{\text{Ber}}^{(i)}$ depends only on $G(X^{(i)})$. Define

$$P^{(i)} = \mathbb{P}\left(G(X^{(i)}) \text{ is connected} \mid Z_{\text{Ber}}^{(i)}\right).$$

Then for any fixed $n \geq 1$, the Rao–Blackwell estimator

$$\widehat{\ell}_{\text{rao}}(\mathcal{G}, p) = \frac{1}{n} \sum_{i=1}^n P^{(i)}$$

is unbiased for $\ell(\mathcal{G}, p)$ and has smaller variance than the crude simulation estimator introduced in (1).

Proof. This proposition follows from standard properties of conditional expectation and the total variance formula. See Billingsley 1995 for further details. \square

Proposition 1 leads to the following algorithm for estimating $\ell(\mathcal{G}, p)$.

Algorithm 1 (Conditional Monte Carlo algorithm for the Bernoulli site percolation model)

1. Set $i = 1$.
2. Generate $N_{\text{deact}}^{(i)} = \min(|\mathcal{V}|, N_{\text{geom}}^{(i)})$, where $N_{\text{geom}}^{(i)}$ has a Geometric (q) distribution on the non-negative integers.
3. If $N_{\text{deact}}^{(i)} = |\mathcal{V}|$, set $P^{(i)} = 1$, set $i = i + 1$ and go to Step 2.
4. Select $N_{\text{deact}}^{(i)}$ vertices uniformly at random from \mathcal{V} without replacement, and denote the chosen vertices by $V_{\text{deact}}^{(i)}$. These vertices will be deactivated.
5. Select a vertex $\omega^{(i)}$ uniformly at random from $\mathcal{V} \setminus V_{\text{deact}}^{(i)}$. This vertex will be activated.
6. Generate the connected component $C_{\text{act}}^{(i)}$ of $G^{(i)}$ containing $\omega^{(i)}$, conditional on the vertices in $V_{\text{deact}}^{(i)}$ being deactivated and $\omega^{(i)}$ being activated.
7. Calculate $P^{(i)} = \mathbb{P}\left(G^{(i)} \text{ is connected} \mid Z_{\text{Ber}}^{(i)} = \left(V_{\text{deact}}^{(i)}, C_{\text{act}}^{(i)}\right)\right)$ according to (3).
8. If $i < n$ set $i = i + 1$ and repeat Step 2. Otherwise return $\frac{1}{n} \sum_{i=1}^n P^{(i)}$.

Note that our construction of Z_{Ber} does not depend on an ordering of the vertices. Another possibility is to take some total ordering of \mathcal{V} and let Z_{Ber} be the connected component of the *first* activated vertex of G . Here ‘first’ is with respect to the ordering of \mathcal{V} . Although we do not pursue this idea further in the discrete case, it leads to a very similar conditional Monte Carlo algorithm to the one described here. We continue this ordering-based approach with reference to random *geometric* graphs in Section 4.

4 CONDITIONAL MONTE CARLO FOR THE GILBERT DISK MODEL

Recall from Section 2 that for the Gilbert disk model the point process ξ is a homogeneous Poisson process with intensity λ in a bounded Borel set R of \mathbb{R}^d . The random graph $G_{\text{geo}} = G_{\text{geo}}(R, \lambda, r)$ is then generated by taking ξ to be the vertices of the graph, and connecting any pair of vertices closer than r in the Euclidean distance by an edge. The probability to be estimated is

$$\ell(R, \lambda, r) = \mathbb{P}(G_{\text{geo}}(R, \lambda, r) \text{ is connected}).$$

Similar to Section 3, we can define the *crude simulation estimator* as

$$\widehat{\ell}_{\text{crude}}(R, \lambda, r) = \frac{1}{n} \sum_{i=1}^n \mathbb{I}\left\{G_{\text{geo}}^{(i)} \text{ is connected}\right\},$$

where $n \geq 1$ is an arbitrary fixed integer and $\left\{G_{\text{geo}}^{(i)}\right\}_{i=1}^{\infty}$ are iid copies of $G_{\text{geo}}(R, \lambda, r)$.

For simplicity we will assume that $d = 2$ and that R is a rectangular region with width w and height h , with bottom left corner at the origin. We will also assume that we have some total ordering of the points of R . Although this ordering can be arbitrary, one natural choice is the *lexicographic ordering*. For $x = (x_1, x_2), y = (y_1, y_2) \in \mathbb{R}^2$, the lexicographic ordering is defined by

$$(x_1, x_2) <_l (y_1, y_2) \quad \text{if and only if} \quad x_1 < y_1 \text{ or } (x_1 = y_1 \text{ and } x_2 < y_2).$$

Another choice is the *distance ordering*, where for some fixed point $z \in R$ the ordering is

$$x <_d y \quad \text{if and only if} \quad \|x - z\| < \|y - z\|.$$

Note that we do not define the ordering among points which are equally distant from z . This is acceptable because we will only apply the ordering to the points of a Poisson process, and with probability 1 there will be no pair of points equally distant from the nonrandom point z .

Let $\eta = (\eta_1, \eta_2)$ be the first point of ξ with respect to the ordering of R . Let Z_{geo} be the vertices of the connected component of G_{geo} that contains η . Then conditional on Z_{geo} there must be no vertices in the region

$$R_{\text{empty}} = \{r \in R \mid r < \eta\} \subseteq R.$$

The set Z_{geo} is equal to $\xi \cap R_{\text{known}}$, where

$$R_{\text{known}} = R \cap \left(\bigcup_{v \in Z_{\text{geo}}} B(v, 2r) \right).$$

On the remainder of R the points of ξ are unknown. That is, conditional on Z_{geo} the distribution of ξ on the region

$$R_{\text{unknown}} = R \setminus (R_{\text{empty}} \cup R_{\text{known}})$$

is that of a homogeneous Poisson point process with intensity λ . The random graph G_{geo} can be connected only if there are no points of ξ in R_{unknown} . Therefore we have

$$\mathbb{P}(G_{\text{geo}} \text{ is connected} \mid Z_{\text{geo}}) = \exp(-\lambda |R_{\text{unknown}}|).$$

The difficulty with applying this formula as part of a conditional Monte Carlo algorithm is determining the area of R_{unknown} , or equivalently $R_{\text{empty}} \cup R_{\text{known}}$. However in some cases this can be relatively straightforward. The following two propositions calculate these areas for the lexicographic ordering and distance ordering.

Proposition 2 (Lexicographic ordering) Consider the *lexicographic ordering* of R . Then

$$R_{\text{empty}} = [0, \eta_1) \times [0, h],$$

and therefore

$$|R_{\text{unknown}}| = h(w - \eta_1) - \left| \left(\bigcup_{v \in Z_{\text{geo}}} B(v, 2r) \right) \cap ([\eta_1, w] \times [0, h]) \right|.$$

Proposition 3 (Distance ordering) Consider the *distance ordering* of R with respect to a fixed point $z \in R$. Then

$$R_{\text{empty}} = B(z, \|z - \eta\|),$$

and therefore

$$|R_{\text{unknown}}| = hw - \left| \left(\bigcup_{v \in Z_{\text{geo}}} B(v, 2r) \cup B(z, \|z - \eta\|) \right) \cap R \right|.$$

Propositions 2 and 3 are easy to prove. The key observation is that we have observed the first point of ξ with respect to the ordering, and this excludes the possibility of observing any other points occurring in a region whose shape depends on the ordering chosen. See Figure 2 for an illustration of the regions R_{known} , R_{unknown} and R_{empty} in the case of the lexicographic ordering.

Applying Propositions 2 and 3, the problem reduces to determining the area of a union of closed balls contained within a rectangular region. In the distance ordering case these balls are of unequal radius. It is well-known (Avis, Bhattacharya, and Imai 1988, Edelsbrunner 1993) that this problem can be solved by constructing the *Laguerre tessellation*, also known as the *power Voronoi diagram*. This diagram can be used to decompose the union of closed balls into a union of simpler regions, each of which is the intersection of a *single* closed ball with finitely many half-planes. If $d = 2$ the diagram can be constructed for n points in $O(n \log n)$ time (Imai, Iri, and Murota 1985). For $d > 2$ the complexity is $O\left(n^{\lfloor \frac{d+1}{2} \rfloor}\right)$ (Aurenhammer 1987).

This leads to the following conditional Monte Carlo estimator for the Gilbert disk model.

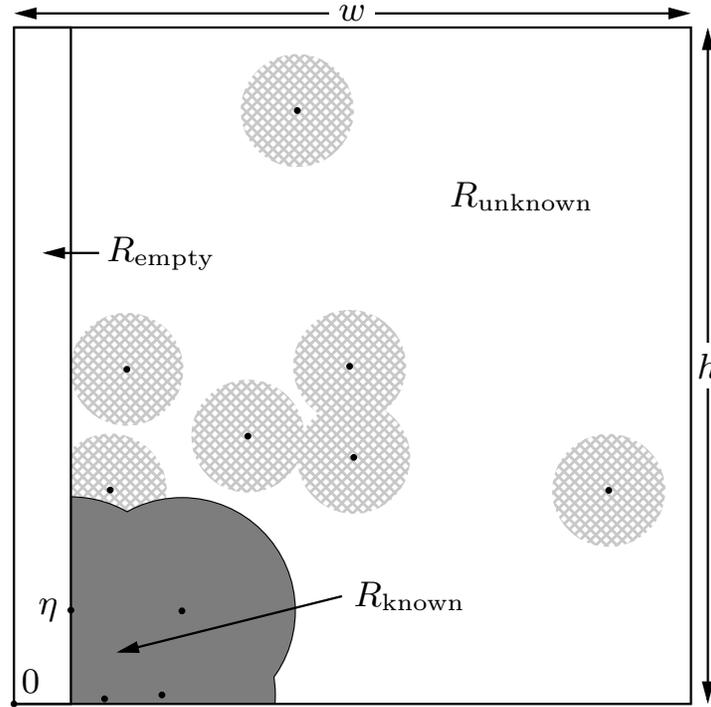


Figure 2: Illustration of the regions R_{known} , R_{unknown} and R_{empty} for the lexicographic ordering on a square region of \mathbb{R}^2 . The region R_{empty} is the rectangular region on the left. The region R_{known} is the region at the bottom left shaded in gray. The region R_{unknown} is the remaining region of R . Crosshatched regions represent other points of ξ that are not in the connected component of η . These regions are included in R_{unknown} .

Proposition 4 (Conditional Monte Carlo estimator for the Gilbert disk model)

Let $\{Z_{\text{geo}}^{(i)}\}_{i=1}^{\infty}$ be iid copies of Z_{geo} and $\{G_{\text{geo}}^{(i)}\}_{i=1}^{\infty}$ be iid copies of $G_{\text{geo}}(R, \lambda, r)$, where each of the $Z_{\text{geo}}^{(i)}$ depend only on $G(X^{(i)})$. Define

$$P_{\text{geom}}^{(i)} = \mathbb{P}\left(G_{\text{geo}}^{(i)} \text{ is connected} \mid Z_{\text{geo}}^{(i)}\right).$$

Then for any fixed $n \geq 1$, the Rao–Blackwell estimator

$$\widehat{\ell}_{\text{rao}}(R, \lambda, r) = \frac{1}{n} \sum_{i=1}^n P_{\text{geo}}^{(i)}$$

is unbiased and has smaller variance than the crude simulation estimator $\widehat{\ell}_{\text{crude}}(R, \lambda, r)$.

For the lexicographic ordering, this leads to the following algorithm. The algorithm for the distance ordering is similar.

Algorithm 2 (Conditional Monte Carlo algorithm for the Gilbert disk model using lexicographic ordering)

1. Set $i = 1$.
2. Generate $Z_{\text{geo}}^{(i)}$ and determine the first point $\eta^{(i)} = (\eta_1^{(i)}, \eta_2^{(i)})$ with respect to the lexicographic ordering.
3. Construct the power Voronoi diagram $V^{(i)}$ of the points in $Z_{\text{geo}}^{(i)}$, with all points having weight $4r^2$.
4. Use $V^{(i)}$ to calculate $|R_{\text{known}}^{(i)}|$.

5. Set $P_{\text{geo}}^{(i)} = \exp\left(-\lambda\left(h\left(w - \eta_1^{(i)}\right) - \left|R_{\text{known}}^{(i)}\right|\right)\right)$.
6. If $i < n$ set $i = i + 1$ and repeat Step 2. Otherwise return $\frac{1}{n} \sum_{i=1}^n P_{\text{geo}}^{(i)}$.

5 NUMERICAL RESULTS

Example 1 Let \mathcal{G} be the 6×6 grid graph without diagonal edges. This graph is small enough to allow complete enumeration of the 2^{36} subgraphs. We can therefore compute the probability of observing a connected subgraph exactly, for any parameter value p . A numerical search for the parameter value which minimized the probability of connectivity gave a value of $p = 0.285$. Note that the whole of \mathcal{G} , any single-vertex subgraph and the empty graph are all connected, so the connectivity probability must approach 1 as p approaches 0 or 1. The probability of connectivity for $p = 0.285$ was calculated to be 0.00125143. For crude Monte Carlo this would result in a relative error of 8.93%.

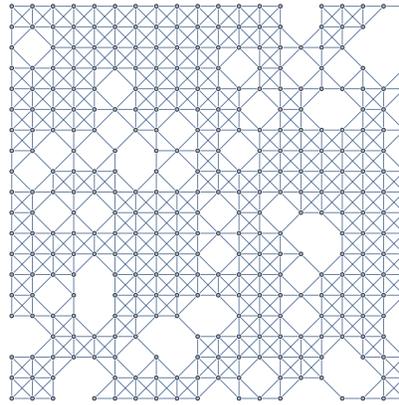
Conditional Monte Carlo was applied with $n = 100,000$ samples, and this was repeated 1,000 times. The average estimate obtained was 1.25×10^{-03} and the estimated relative error was 2.48%.

Example 2 We started with a 20×20 grid graph which included diagonal edges and generated a random subgraph by retaining at random 340 of the 400 vertices. The base graph that was generated is shown in Figure 3. In this case exact computation is difficult, and we actually simulate the crude Monte Carlo method. Both crude Monte Carlo and conditional Monte Carlo were applied for 20 different parameter values between $p = 0.05$ and $p = 0.99$. For values of p between 0.1 and 0.33 inclusive the crude method did not identify any connected subgraphs and therefore estimated a probability of 0. A sample of the results where both methods estimated non-zero probabilities is shown in Table 1. An up to five-fold improvement in relative error is observed when using the Rao–Blackwell estimator as compared to the crude estimator.

Table 1: Simulation results for a randomly generated subgraph of the 20×20 grid graph.

Method	p	Estimate	Relative Error %
Crude	0.05	5.98×10^{-07}	130
Crude	0.43	1.56×10^{-07}	251
Crude	0.47	1.17×10^{-05}	30
Crude	0.52	2.99×10^{-04}	5.8
Crude	0.57	3.33×10^{-03}	1.7
Conditional	0.05	5.72×10^{-07}	25
Conditional	0.43	1.75×10^{-07}	68
Conditional	0.47	1.16×10^{-05}	11.2
Conditional	0.52	3.00×10^{-04}	3.1
Conditional	0.57	3.34×10^{-03}	1.0

Example 3 We considered the Gilbert disk model on a 6 by 6 square region of \mathbb{R}^2 . The homogeneous Poisson point process generating the vertices of the graph had intensity 10, and the distance r at which points are connected was allowed to be 0.20 or 0.21. Both the distance and lexicographic orderings were considered. In the case of the distance ordering the fixed point z was taken to be the center of R . We used $n = 10,000$ samples for all three estimators, and this was repeated 1,000 times to estimate the relative error. The simulation results are shown in Table 2. The distance ordering appears to outperform the lexicographic ordering, and gives an up to 7-fold improvement in the relative error compared to the crude simulation estimator.

Figure 3: Subgraph of the 20×20 grid graph used as the base graph in Example 2.Table 2: Simulation results for the Gilbert disk model on a 6×6 region with intensity 10.

Method	r	Estimate	Relative Error %
Crude	0.20	4.90×10^{-05}	463
Crude	0.21	1.34×10^{-03}	86
Lexicographic	0.20	5.11×10^{-05}	52
Lexicographic	0.21	1.34×10^{-03}	12
Distance	0.20	4.90×10^{-05}	46
Distance	0.21	1.34×10^{-03}	12

ACKNOWLEDGMENTS

This work was supported by the ARC Center of Excellence in Mathematical and Statistical Frontiers for Big Data, Big Models and New Insights (ACEMS, CE140100049), and by the DAAD / Go8 Australia – Germany Joint Research Cooperation Scheme.

REFERENCES

- Asmussen, S., and P. W. Glynn. 2007. *Stochastic Simulation : Algorithms and Analysis*. New York: Springer.
- Asmussen, S., and R. Y. Rubinstein. 1995. “Steady state rare events simulation in queueing models and its complexity properties”. In *Advances in Queueing*, 429–461. Boca Raton: CRC Press.
- Aurenhammer, F. 1987. “Power diagrams: Properties, algorithms and applications”. *SIAM J. Comput.* 16 (1): 78–96.
- Avis, D., B. Bhattacharya, and H. Imai. 1988. “Computing the volume of the union of spheres”. *The Visual Computer* 3 (6): 323–328.
- Billingsley, P. 1995.. *Probability and Measure*. 3rd ed. New York: J. Wiley & Sons.
- Bollobás, B., and O. Riordan. 2006. *Percolation*. Cambridge: Cambridge University Press.
- Botev, Z. I., and D. P. Kroese. 2012. “Efficient Monte Carlo simulation via the generalized splitting method”. *Statistics and Computing* 22 (1): 1–16.
- Brereton, T., C. Hirsch, D. P. Kroese, and V. Schmidt. 2014. “Pair connectedness and shortest-path scaling in critical continuum percolation”. Submitted.
- Chiu, S. N., D. Stoyan, W. S. Kendall, and J. Mecke. 2013. *Stochastic Geometry and Its Applications*. Chichester: J. Wiley & Sons.
- Colbourn, C. J. 1987. *The Combinatorics of Network Reliability*. New York: Oxford University Press, Inc.

- Edelsbrunner, H. 1993. “The union of balls and its dual shape”. In *Proceedings of the Ninth Annual Symposium on Computational Geometry*, SCG '93, 218–231: ACM.
- Garvels, M. J. J., J.-K. C. W. van Ommeren, and D. P. Kroese. 2002. “On the importance function in splitting simulation”. *European Transactions on Telecommunications* 13 (4): 363–371.
- Gertsbakh, I. B., and Y. Shpungin. 2010. *Models of Network Reliability*. Boca Raton: CRC Press.
- Glasserman, P., P. Heidelberger, P. Shahabuddin, and T. Zajic. 1999. “Multilevel splitting for estimating rare event probabilities”. *Operations Research* 47 (4): 585–600.
- Glynn, P. W., and D. L. Iglehart. 1989. “Importance sampling for stochastic simulations”. *Management Science* 35 (11): 1367–1392.
- Imai, H., M. Iri, and K. Murota. 1985. “Voronoi diagram in the Laguerre geometry and its applications”. *SIAM Journal on Computing* 14 (1): 93–105.
- Kahn, H., and T. E. Harris. 1951. “Estimation of particle transmission by random sampling”. *National Bureau of Standards Applied Mathematics Series* 12:27–30.
- Kroese, D. P., T. Taimre, and Z. I. Botev. 2011. *Handbook of Monte Carlo Methods*. New York: J. Wiley & Sons.
- L’Ecuyer, P., V. Demers, and B. Tuffin. 2006. “Splitting for rare-event simulation”. In *Proceedings of the 38th Conference on Winter Simulation*, WSC '06, 137–148.
- Li, J., and M. Östling. 2013. “Percolation thresholds of two-dimensional continuum systems of rectangles”. *Phys. Rev. E* 88:012101.
- Quintanilla, J. A., and R. M. Ziff. 2007. “Asymmetry in the percolation thresholds of fully penetrable disks with two different radii”. *Phys. Rev. E* 76:051115.
- Rubinstein, R. Y., and D. P. Kroese. 2008. *Simulation and the Monte Carlo Method*. 2nd ed. New York: J. Wiley & Sons.
- Sahini, M., and M. Sahimi. 1994. *Applications of Percolation Theory*. Bristol: Taylor & Francis.
- Stenzel, O., L. J. A. Koster, R. Thiedmann, S. D. Oosterhout, R. A. J. Janssen, and V. Schmidt. 2012. “A new approach to model-based simulation of disordered polymer blend solar cells”. *Advanced Functional Materials* 22 (6): 1236–1244.
- Torquato, S., and Y. Jiao. 2012. “Effect of dimensionality on the continuum percolation of overlapping hyperspheres and hypercubes. II. Simulation results and analyses”. *The Journal of Chemical Physics* 137 (7): 074106.
- Weichenberg, G., V. W. S. Chan, and M. Medard. 2004. “High-reliability topological architectures for networks under stress”. *IEEE Journal on Selected Areas in Communications* 22 (9): 1830–1845.

AUTHOR BIOGRAPHIES

ROHAN SHAH is a PhD student at the School of Mathematics and Physics of the University of Queensland. He has an honors degree in statistics from the University of Western Australia. His research interests include Monte Carlo methods, rare event simulation, stochastic geometry and statistical software. His email address is rohan.shah@uqconnect.edu.au.

CHRISTIAN HIRSCH is a PhD student at the Faculty of Mathematics and Economics of Ulm University. He has a Master (Diploma) in Mathematics from the Ludwig Maximilian University of Munich. His research interests include stochastic geometry, random geometric graphs, and Monte Carlo simulation of spatial stochastic models. His personal website can be found under <http://www.uni-ulm.de/stochastik>. His email address is christian.hirsch@uni-ulm.de.

DIRK P KROESE is a professor of Mathematics and Statistics at the University of Queensland. He is the co-author of several influential monographs on simulation and Monte Carlo methods, including *Handbook of Monte Carlo Methods and Simulation and the Monte Carlo Method*, (2nd Edition). Dirk is a pioneer of the well-known Cross-Entropy method — an adaptive Monte Carlo technique, invented by Reuven Rubinstein,

which is being used around the world to help solve difficult estimation and optimization problems in science, engineering, and finance. His personal website can be found under <http://www.maths.uq.edu.au/~kroese>. His email address is kroese@maths.uq.edu.au.

VOLKER SCHMIDT is Professor at the Faculty of Mathematics and Economics of Ulm University. His research interests include stochastic geometry, spatial statistics, and Monte Carlo simulation of spatial stochastic models as well as their applications to structural analysis of (microscopic and geographically mapped) image data. He is (co-) author of more than 100 peer-reviewed publications, including several textbooks and monographs. His personal website can be found under <http://www.uni-ulm.de/stochastik>. His email address is volker.schmidt@uni-ulm.de.