

# 1 Subset selection via continuous optimization with 2 applications to network design

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6 **Abstract** Choosing a subset of representative items from a set of alternatives  
7 is an important problem in many scientific fields such as environmental science  
8 and statistics. For most practical problems, however, a computationally efficient  
9 solution method is not known to exist. While this problem has attracted a signif-  
10 icant amount of attention, the majority of specifically designed algorithms do not  
11 scale well with respect to the problem size or do not provide a usable open-source  
12 package. In this study, we show that any global continuous optimization technique  
13 can be used for solving the representative subset selection problem. The latter is  
14 achieved by designing a simple transformation which embeds the problem's dis-  
15 crete space into a larger continuous space. The proposed methodology is applied  
16 to design problems in environmental and statistical domains. We evaluate the pro-  
17 posed method using several open-source global optimization packages, and show  
18 that this technique compares favorably with existing direct methods.

19 **Keywords** Ozone · Monitoring network design · D-optimal experimental design ·  
20 global optimization · space embedding

## 21 1 Introduction

22 Networks are pervasive in modern society. Environmental and distribution net-  
23 works (waste-water, power grids, aviation, World Wide Web), Social networks  
24 (Facebook, YouTube, Twitter), and biological networks (neural, metabolic, protein-  
25 protein interaction) all play a key role in the functioning of our world. Monitoring  
26 such networks is now an important part of scientific endeavor. For example, an im-  
27 proper monitoring of waste-water or carbon dioxide emission can result in a serious  
28 damage to the ecosystem (Stokes and Horvath, 2010; Mi et al., 2017; Park et al.,  
29 2013; Chan and Yao, 2008; Ramanathan and Carmichael, 2008; Bruns et al., 1991;

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30 Wiersma, 1984; von Brömssen et al., 2018). Likewise, a diffusion of false informa-  
31 tion from social media (“fake news”), can potentially have a devastating effect on  
32 national security and political stability (Allcott and Gentzkow, 2017; Silverman  
33 and Singer-Vine, 2016). In this study, we address the *monitoring network design*  
34 problem (Le and Zidek, 2006) by introducing a general technique that allows one  
35 to use *any* global continuous optimization method for solving the problem. The  
36 proposed method provides high-quality solutions to the monitoring network design  
37 problem with minimal development effort.

38 Regardless of the application domain, it is convenient to define a network as a  
39 graph with  $m$  vertices. These vertices can represent, for example, possible locations  
40 for a placement of wastewater monitoring sensor, web-servers, influential bloggers,  
41 or electrical re-transmission blocks. Due to the high operational cost associated  
42 with the size of modern networks, monitoring the entire system is generally un-  
43 feasible. Consequently, a natural approach is to analyze some representative part  
44 of the network in order to infer the entire network state. In this way, the optimal  
45 monitoring network design task can be viewed as an *optimal fixed-size subset se-*  
46 *lection* (OFSS) problem in which a subset of size  $k \ll m$  vertices (from  $m$  initial  
47 vertices) has to be chosen such that a certain *utility function* is maximized. Under  
48 the monitoring network design setting, the utility function will generally measure  
49 the information gain obtained from a decision of choosing a particular subset of  
50 vertices.

51 Because of the importance of the OFSS problem from both a theoretical and  
52 practical point of view, it has attracted a significant amount of research atten-  
53 tion (Wolters, 2015; Chao et al., 2015; Yu and Yuan, 1992). However, excluding  
54 some trivial cases, this problem belongs to the NP-hard complexity class (Nataraj-  
55 an, 1995; Geoffrey et al., 1997; Chun-Wa et al., 1995), necessitating approximate  
56 solution methods. Due to the hardness result, no approximation algorithm can  
57 guarantee the discovery of the optimal solution to the OFSS problem. From the  
58 practical point of view, however, it is beneficial to implement a number of different  
59 approximation methods. In this way, the designer can then solve the problem with  
60 several algorithms, and adopt the best solution found.

61 Despite the problem importance, there exists a deficit of freely available and ac-  
62 cessible software that is specifically designed to handle the OFSS problem (Ramiro  
63 et al., 2010; Wolters, 2015). In fact, to the best of our knowledge, the only open-  
64 source package available is `kofnGA` (Wolters, 2015). This study aims to address the  
65 above gap by introducing a simple transformation, which embeds the (discrete)  
66 solution space of the OFSS problem into a larger continuous space, allowing the  
67 practitioner to use any freely available or proprietary continuous global optimiza-  
68 tion software to obtain a solution to the OFSS problem. Since the majority of  
69 global optimization packages (Mullen, 2014) use different heuristics, the proposed  
70 technique introduces a significant practical advantage, in the sense that one can  
71 take advantage of diverse algorithms with minimal development effort.

72 The rest of the study is organized as follows. In Section 2 we formulate the  
73 OFSS problem and show that both the optimal network monitoring design and the  
74 D-optimal experimental design problems fit into the OFSS framework. Then, we  
75 give a brief overview of the modern continuous global optimization research, and  
76 introduce a simple transformation from the continuous to the original discrete state  
77 space. In Section 3 we perform a detailed experimental study using several popular

78 open-source global optimization packages. Finally, in Section 4 we summarize our  
 79 findings and discuss possible directions for future research.

## 80 2 Methodology

### 81 2.1 The OFSS problem

The formal definition of the OFSS problem is as follows. Let  $\mathcal{L} = \{1, \dots, m\}$  be the set of indices that corresponds, for example, to  $m$  available locations for a placement of wastewater monitoring sensor. Since we wish to select  $k \leq m$  indices from  $\mathcal{L}$ , the set of all possible solutions of the  $k$ -sized OFSS problem is defined via

$$\mathcal{X} = \left\{ \mathbf{x} \in \mathcal{P}(\{1, \dots, m\}), |\mathbf{x}| = k \right\}, \quad (1)$$

82 where  $\mathcal{P}$  stands for the power-set, and  $|\mathbf{x}|$  denotes the cardinality of  $\mathbf{x}$ . Recall  
 83 that the selection should be performed subject to a maximization of certain utility  
 84 criterion. Therefore, we define a *utility function*  $U : \mathcal{X} \rightarrow \mathbb{R}$ , which we seek to  
 85 maximize. Finally, the OFSS problem can be formulated via:

$$\max_{\mathbf{x} \in \mathcal{X}} U(\mathbf{x}), \quad (2)$$

86 while the set of optimal indices is available from the solution of  $\operatorname{argmax}_{\mathbf{x} \in \mathcal{X}} U(\mathbf{x})$ .  
 87 In this study, we consider two real-life applications that fall into the OFSS problem  
 88 setting; the optimal monitoring network design, and the D-optimal experimental  
 89 design.

#### 90 *Entropy-based environmental monitoring network design*

91 Environmental monitoring networks design involves the selection of  $k$  representa-  
 92 tive locations out of  $m \geq k$  possible locations in a region of interest. For example,  
 93 Fig. 1 from (Wolters, 2015) shows  $m = 100$  potential locations (left panel), and  
 94  $k = 9$  actual locations (right panel) of ozone monitoring stations in the state of  
 95 New York (Le and Zidek, 2006).

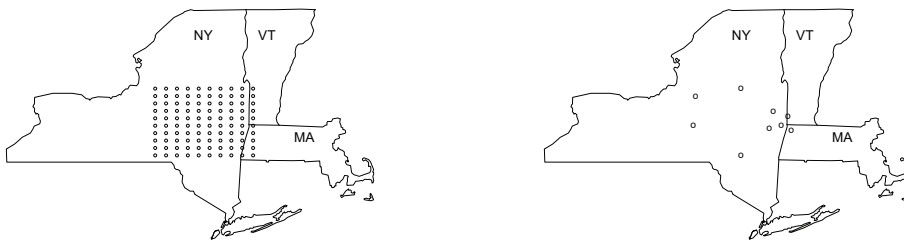


Fig. 1: Potential (left panel), and existing (right panel), ozone monitoring sites in the state of New York.

96 Since the cost of auditing of all possible locations can be prohibitively high, we  
 97 generally wish to select  $k \ll m$  representative sites for the placement of monitoring  
 98 stations. The selection is performed with a view to maximize some utility function  
 99 such as an obtained information gain from the selection of a particular set of  $k$   
 100 monitoring sites. Formally, the ozone concentration data that is collected from a  
 101 certain site can be modeled via a stochastic time series. In addition, it is convenient  
 102 to model the data collected from all sites as a random field  $\{Z(l) : l \in D \subset \mathbb{R}^d\}$ ,  
 103 where  $D$  generally represents a discretized subset of  $\mathbb{R}^d$ . Our task is to select  $k$   
 104 locations from a finite set  $D$ . The monitoring network design problem is usually  
 105 associated with the two-dimensional or the three-dimensional space, that is, we  
 106 generally take  $d = 2$  or  $d = 3$ .

107 To put the monitoring networks design problem into the OFSS setting, we need  
 108 to specify the utility function. Note that the  $m$ -dimensional vector of observations  
 109  $\mathbf{Z}$ , which is obtainable from the entire set of  $m$  sites, can be partitioned into  
 110 two vectors  $\mathbf{Z}^{(1)}$  and  $\mathbf{Z}^{(2)}$ , where  $\mathbf{Z}^{(1)}$  represents  $m - k$  observations from the  
 111 set of *unmeasured* sites and  $\mathbf{Z}^{(2)}$  represents the remaining  $k$  observations from  
 112 the set of *measured* sites. Let  $\mathcal{D}(\mathbf{z})$  be the joint probability density of  $\mathbf{Z}$ . Then,  
 113 the total uncertainty about  $\mathbf{Z}$  is given by the entropy  $H(\mathbf{Z}) = \mathbb{E}[-\ln \mathcal{D}(\mathbf{Z})] =$   
 114  $\mathbb{E}[-\ln \mathcal{D}(\mathbf{Z}^{(1)}, \mathbf{Z}^{(2)})]$  (Shannon, 1948). Under the above setting, only the  $\mathbf{Z}^{(2)}$   
 115 vector is observable, so we are interested in minimizing the uncertainty about  $\mathbf{Z}^{(1)}$   
 116 given  $\mathbf{Z}^{(2)}$ . That is, we seek to minimize the conditional entropy  $H(\mathbf{Z}^{(1)} | \mathbf{Z}^{(2)})$ .  
 117 From the chain rule of conditional entropy, namely, from

$$H(\mathbf{Z}^{(1)} | \mathbf{Z}^{(2)}) = H(\mathbf{Z}^{(1)}, \mathbf{Z}^{(2)}) - H(\mathbf{Z}^{(2)}),$$

118 combined with the fact that the total system entropy  $H(\mathbf{Z}) = H(\mathbf{Z}^{(1)}, \mathbf{Z}^{(2)})$   
 119 is fixed, we conclude that the minimization of  $H(\mathbf{Z}^{(1)} | \mathbf{Z}^{(2)})$  is equivalent to  
 120 the maximization of  $H(\mathbf{Z}^{(2)})$ . Namely, our task is to maximize the entropy of  
 121 observations associated with the observed sites.

122 The definition of the corresponding utility function  $U : \mathcal{X} \rightarrow \mathbb{R}$  is straightfor-  
 123 ward. Specifically, for  $\mathbf{x} = \{x_1, \dots, x_k\} \in \mathcal{X}$ , let  $\mathbf{Z}_{\mathbf{x}} = (Z_{x_1}, \dots, Z_{x_k})$  be a  $k$ -sized  
 124 vector of measurements which was extracted from  $\mathbf{Z} = (Z_1, \dots, Z_m)$  using the set  
 125 of indices  $\mathbf{x}$ . Then, by defining  $U(\mathbf{x}) \triangleq H(\mathbf{Z}_{\mathbf{x}})$ , we establish the correspondence  
 126 between the optimal network design and the OFSS problems.

127 While the above setup is valid for a general joint probability density  $\mathcal{D}(\mathbf{z})$ ,  
 128 we concentrate on the important special case, for which  $\mathcal{D}(\mathbf{z})$  is the multivariate  
 129 Normal distribution; namely,  $\mathbf{Z} \sim \mathbf{N}_m(\boldsymbol{\mu}, \boldsymbol{\Sigma})$ . For this particular distribution, we  
 130 have that  $H(\mathbf{Z}) \propto \ln |\boldsymbol{\Sigma}|$ , where  $|\boldsymbol{\Sigma}|$  is the determinant of the covariance matrix  
 131 (Ahmed and Gokhale, 2006). Therefore, the maximization of  $H(\mathbf{Z}^{(2)})$  requires  
 132 the maximization of the natural logarithm of a *principal sub-matrix* of  $\boldsymbol{\Sigma}$ , indexed  
 133 by the observable sites index set  $\mathbf{x}$ . Such principal sub-matrix is denoted by  $\boldsymbol{\Sigma}_{\mathbf{x}}$   
 134 and one can extract it from  $\boldsymbol{\Sigma}$  by taking all rows and columns of  $\boldsymbol{\Sigma}$  that are in  $\mathbf{x}$ .  
 135 Finally, we define the utility function to be  $U(\mathbf{x}) \triangleq \ln |\boldsymbol{\Sigma}_{\mathbf{x}}|$ .

136 *D-optimal experimental design*

137 The D-optimal experimental design problem involves a discovery of the best subset  
 138 of  $m$  possible experiments. Formally, let  $X$  be an  $m \times m$  model matrix. The D-  
 139 optimal design is one that maximizes the determinant of the principal sub-matrix,  
 140  $X_{\mathbf{x}}^{\top} X_{\mathbf{x}}$ , where  $X_{\mathbf{x}}$  is a  $k \times k$  matrix ( $k \leq m$ ) extracted from  $X$  by taking all rows  
 141 and columns of  $X$  that are indexed by  $\mathbf{x} \in \mathcal{X}$  (Goos and Bradley, 2011). The  
 142 objective is to maximize the so-called D-optimality criterion. The latter is defined  
 143 as a negative of natural logarithm of the experiment's *information matrix*  $X_{\mathbf{x}}^{\top} X_{\mathbf{x}}$   
 144 (Roy, 2000). Therefore, by defining the utility function to be  $U(\mathbf{x}) \triangleq \ln |X_{\mathbf{x}}^{\top} X_{\mathbf{x}}|$ ,  
 145 we can easily see that the D-optimal experimental design problem corresponds to  
 146 the OFSS setting (Wolters, 2015).

147 Our objective is to develop an efficient procedure for the maximization of (2).  
 148 However, since this problem is known to be hard (Chun-Wa et al., 1995), we  
 149 resort to approximate evolutionary strategies. A brief overview of these methods  
 150 is detailed in Section 2.2. We refer to (Mullen, 2014) for a detailed discussion of  
 151 global optimization techniques.

## 152 2.2 Global continuous optimization

153 Evolutionary strategies are stochastic optimiza-  
 154 tion procedures that are motivated by a biolog-  
 155 ical process of natural selection (Bäck, 1996).  
 156 Specifically, these methods simulate the evolu-  
 157 tionary mechanism by producing a sequence of  
 158 population generations that (hopefully) improve  
 159 the population's average fitness (utility) over  
 160 time. A very general evolutionary framework,  
 161 which is implemented in many optimization  
 162 packages, is summarized in Fig. 2. In this study,  
 163 we consider several state-of-the-art evolutionary  
 164 algorithms, namely, Differential Evolution (DE)  
 165 (Price et al., 2005), Particle Swarm Optimization  
 166 (PSO) (Kennedy and Mendes, 2002), Genetic al-  
 167 gorithm (GA) (Melanie, 1998), and Cross En-  
 168 tropy algorithm (CE) (Rubinstein and Kroese,  
 169 2017).

170 The major advantage of such nature-inspired  
 171 methods, is that they can operate in both the  
 172 discrete and the continuous state spaces. How-  
 173 ever, for non-standard discrete applications such  
 174 as the OFSS problem, an evolutionary method  
 175 will customarily require an adjustment which  
 176 will generally involve a development of a spe-  
 177 cialized *next population creation* (see Fig. 2). De-  
 178 signing such operations can be a time-consuming  
 179 task that may require a substantial development  
 180 effort. For example, the **kofnGA** (Wolters, 2015)

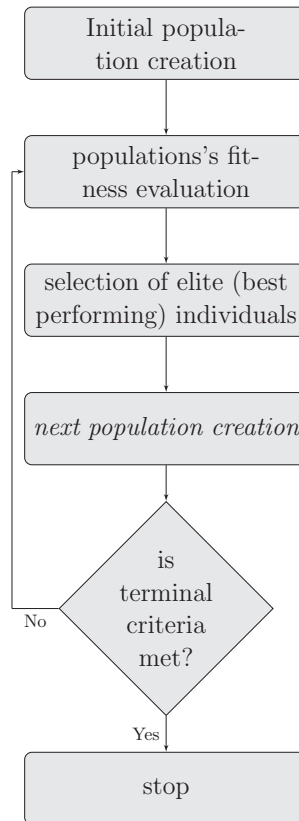


Fig. 2: A general evolutionary framework.

181 package, which is suitable for handling OFSS tasks, such as the optimal monitoring  
 182 network design problems, is based on the discrete version of GA (Melanie, 1998).

183 Motivated by the fact that the next population creation in the continuous  
 184 space is quite standard (Mullen, 2014), we explore a possibility of adopting the  
 185 continuous optimization approach. In particular, we design a *transformation* that  
 186 allows us to use an embedding of the original discrete space of the OFSS problem  
 187 into a larger continuous space. The transformation is detailed in Section 2.3.

### 188 2.3 The transformation

In this section, we assume that one has an access to a global continuous optimization procedure (such as DE, PSO, GA, or CE), which is capable of handling problems of the type:

$$\begin{aligned} & \max_{\mathbf{y} \in \mathbb{R}^m} C(\mathbf{y}) & (3) \\ & \text{subject to: } \mathbf{y} \in \mathcal{Y}, \end{aligned}$$

189 where  $C : \mathbb{R}^m \rightarrow \mathbb{R}$  is an objective function, and  $\mathcal{Y} \subseteq \mathbb{R}^m$  is a compact convex set.  
 190 The transformation from the continuous state space  $\mathcal{Y}$  to the original discrete state  
 191 space  $\mathcal{X}$  (see (1)), and the corresponding calculation of the utility, is summarized  
 192 in Algorithm 1.

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#### Algorithm 1: The transformation

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**Input:** A vector of real numbers  $\mathbf{y} \in \mathcal{Y} \subseteq \mathbb{R}^m$ ,  $m, k \in \mathbb{N}$  such that ( $k \leq m$ ), and a utility function  $U : \mathcal{X} \rightarrow \mathbb{R}$  from (2).

**Output:** The utility of  $\mathbf{x} \in \mathcal{X}$ , where  $\mathbf{x}$  corresponds to the continuous input  $\mathbf{y} \in \mathcal{Y}$ .

```

1  $\mathbf{x}' \triangleq (x_1, \dots, x_m) \leftarrow \text{argsort}(\mathbf{y})$ 
2  $\mathbf{x} \leftarrow \{x_1, \dots, x_k\}$ 
3 return  $U(\mathbf{x})$ 

```

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193 In Line 1 of Algorithm 1, an *ascending* sorting procedure (Cormen et al., 2001)  
 194 is used to determine the internal ordering of indices of  $\mathbf{y}$ . Note that  $\mathbf{x}'$  is a permutation  
 195 of  $(1, \dots, m)$ . Then, in line 2, we fix  $\mathbf{x}$  to be a set that contains the  $k$ -sized  
 196 prefix of  $\mathbf{x}'$ ; note that  $\mathbf{x} \in \mathcal{X}$ . Finally, in Line 3 we use the previously obtained  $\mathbf{x}$   
 197 to calculate the utility in the discrete space.

198 It is not very hard to see that there exists  $\mathbf{y}^* \in \mathcal{Y}$  that corresponds (via  
 199 the transformation in Algorithm 1), to the optimal (discrete) solution  $\mathbf{x}^* \in \mathcal{X}$ ,  
 200 such that  $U(\mathbf{x}^*) \geq U(\mathbf{x})$  for all  $\mathbf{x} \in \mathcal{X}$ . To see this, suppose without the loss of  
 201 generality that  $\mathbf{x}^*$  contains the first  $k$  sites, that is,  $\mathbf{x}^* = \{1, \dots, k\}$ . Then, for  
 202 example,  $\mathbf{y}^* = (\underbrace{0, \dots, 0}_k, \underbrace{1, \dots, 1}_{n-k})$  corresponds to  $\mathbf{x}^* = \{1, \dots, k\}$  via Algorithm 1.

203 The proposed transformation opens the way for addressing the OFSS problem  
 204 via any global continuous optimization method. To see this, note that Algorithm 1  
 205 implements the function  $C(\mathbf{y})$  in (3). In Section 3 we perform a comprehensive  
 206 experimental study to examine the performance of the proposed transformation.

### 3 Results and discussion

We investigate the accuracy of the proposed method when applied to several representative examples. The rest of the section is structured as follows.

1. The first example is a custom-made environmental monitoring design task for which we know the optimal solution. Specifically, we examine two regular lattices of dimensions  $5 \times 5$  and  $9 \times 9$ . To ensure symmetry, we choose  $k = 9$  and  $k = 25$  (see Fig. 3), for the first and the second lattice, respectively. This example enables the precise bench-marking of the proposed algorithm accuracy.
2. Our second problem is a real-life example in which we consider the ozone concentrations in the state of New York. The example involves an analysis of 100 potential locations for a placement of 15 monitoring stations.
3. Next, we examine a larger monitoring network design instance with 1000 possible site placements among which we seek to determine the optimal location of 10, 15, and 20 monitoring stations. This example is motivated by the fact that nowadays, a monitoring network designer encounters typical networks with several hundreds of potential locations.
4. To demonstrate the generality of the proposed technique, we conclude the numerical study with an example from a quite different domain. Specifically, we consider the optimal design of statistical experiments and in particular the D-optimal design problem (Goos and Bradley, 2011). The application of the method to the D-optimal design problem is feasible, since the latter can be modeled via the OFSS setting.

#### *The experimental setup*

All tests were implemented using the open-source R statistical software package (version 3.5.2). The software is available from the author's web-page along with all examples from this section. The script was executed on 64 bit Windows 10 desktop machine with Intel Core i7-3770 quad-core 3.4GHz processor and 16GB of RAM.

With a view to provide a fair comparison between different global optimization algorithms, we use a set of freely available R packages. Specifically, we work with the following libraries: `kofnGA` version 1.3, `DEoptim` version 2.2.4, `pso` version 1.0.3, `GA` version 3.1.1, and `CEoptim` version 1.2. In addition, we ran a number of preliminary benchmarks (not reported here), to determine a reasonably robust parameter setting for the optimization packages under consideration. While the optimal parameter determination is subject to a further research, our objective was to allow an approximately equal running (CPU) time for all methods involved.

Finally, note that we seek to *maximize* the entropy of the measurements obtained from the set of observable sites. In this section, we rather minimize the corresponding negative entropy. The latter is motivated by a purely technical consideration, since the majority of global optimization packages minimize the objective function. Specifically, instead of (3), we work with:

$$\min_{\mathbf{y} \in \mathbb{R}^m} -C(\mathbf{y}), \quad \text{subject to: } \mathbf{y} \in \mathcal{Y}.$$

Unless stated otherwise, each algorithm (`kofnGA`, `DE`, `PSO`, `GA`, and `CE`), was executed for 20 times. All CPU times are reported in seconds. To ensure repro-

250 ducibility, we use the same random seed (of “12345”) for all experiments reported  
 251 in this section.

### 252 3.1 Custom grid with known solution

253 In order to benchmark the proposed method, we consider a custom 2D-grid (reg-  
 254 ular lattice) of dimensions  $5 \times 5$  and  $9 \times 9$ . Note that these examples represent  
 255 random fields with  $m = 25$  and  $m = 81$  possible locations, respectively. In these  
 256 benchmarks, we assume that the data is distributed according to the multivariate  
 257 Normal distribution, namely  $\mathbf{Z} \sim \mathbf{N}_m(\boldsymbol{\mu}, \boldsymbol{\Sigma})$ , and that the covariance matrix  $\boldsymbol{\Sigma}$   
 258 is known. Specifically, for each component  $\Sigma_{i,j}$  for  $1 \leq i, j \leq m$  in  $\boldsymbol{\Sigma}$ , we set  
 259  $\Sigma_{i,j} = 10 \exp\{-0.1 \text{dist}(i, j)\}$ , where  $\text{dist}(i, j)$  is the Euclidean distance between  
 260 locations  $i$  and  $j$  in the lattice. We set  $k = 9$  and  $k = 25$  for the  $5 \times 5$  and the  
 261  $9 \times 9$  grids, respectively. The corresponding optimal locations of the monitoring  
 stations are shown in Fig. 3.

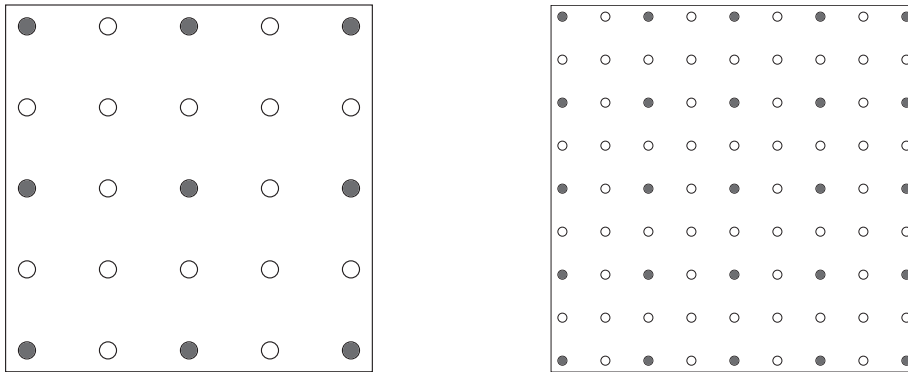


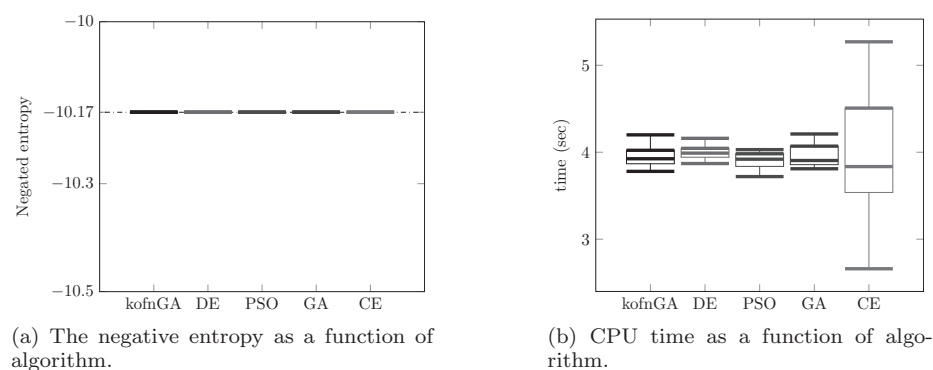
Fig. 3: Optimal locations for the  $5 \times 5$   $k = 9$  (left panel), and  $9 \times 9$   $k = 25$  (right panel), lattices.

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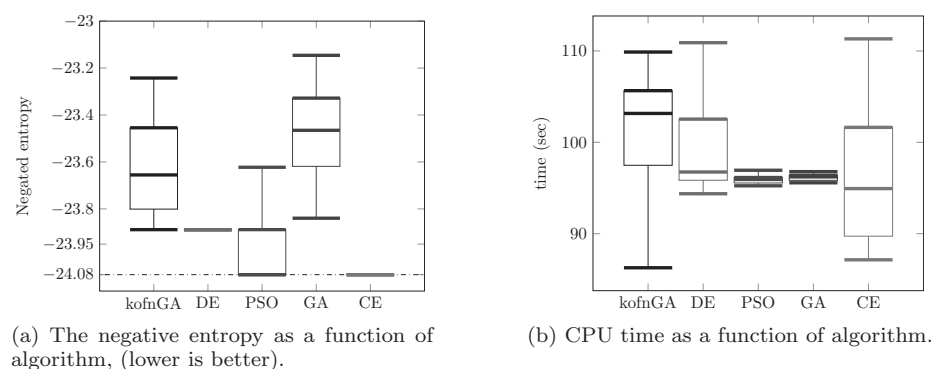
#### 263 *The $5 \times 5$ lattice*

264 We start with a small  $5 \times 5$  lattice for which we wish to select  $k = 9$  observable  
 265 sites from  $m = 25$  possible locations. In this particular setting, a full-enumeration  
 266 procedure requires to consider  $\binom{25}{9} = 2042975$  combinations. The optimal objec-  
 267 tive function value found is  $-10.167694$  and the obtained optimization results  
 268 are shown in Fig. 4. As we are dealing with a small toy example, it is not very  
 269 surprising to see that all methods deliver the optimal solution.



Fig. 4: The  $5 \times 5$  lattice optimization results.270 *The  $9 \times 9$  lattice*

271 Next, we continue with a bigger  $9 \times 9$  lattice for which we set  $m = 81$  and  $k = 25$ . In  
 272 this situation, a full-enumeration procedure needs to consider  $\binom{81}{25} \approx 5.256 \times 10^{20}$   
 273 combinations. The optimal objective function value found is  $-24.08018$ . Fig. 5  
 274 summarizes the obtained optimization results. For this particular problem, CE  
 275 always found the optimal solution, and the PSO was next to the best. Similar to  
 276 the  $5 \times 5$  lattice example, the CE running time introduced a higher variance.

Fig. 5: The  $9 \times 9$  lattice results obtained with kofnGA, DE, PSO, GA, and CE.277 **3.2 Environmental monitoring**

278 In this section, we consider a real-life network monitoring design problem from (Le  
 279 and Zidek, 2006). Since the authors examine 100 possible locations for placing 15

280 monitoring stations, the full-enumeration procedure requires the consideration of  
 281  $\binom{100}{15} \approx 2.533 \times 10^{17}$  combinations.

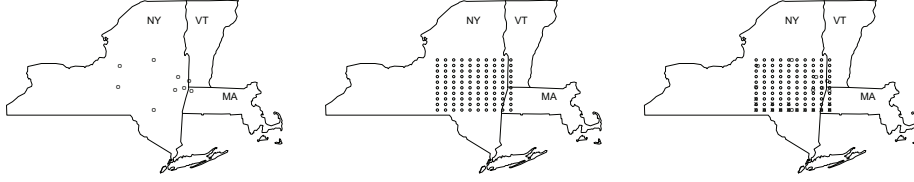
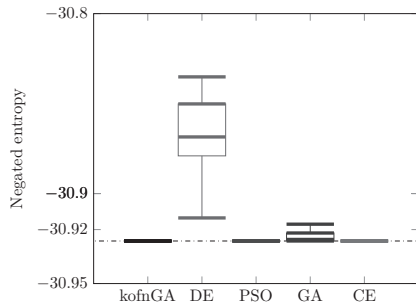


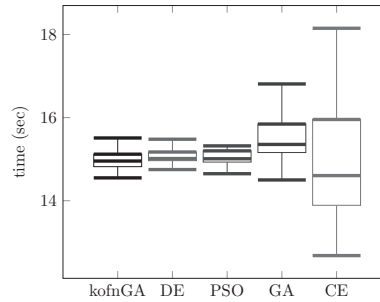
Fig. 6: Existing sites (left panel), potential placements (middle panel), and a combination of existing ( $\circ$ ) and optimal new ( $\otimes$ ) site placement (right panel).

282 The left and the middle panel of Fig. 6 show the 9 existing and the 100 potential  
 283 site locations, respectively. The Fig. 6 right panel shows the nine original ( $\circ$ )  
 284 locations and the new placement of optimized locations ( $\otimes$ ) of 15 monitoring  
 285 stations.

286 Fig. 7 summarizes the optimization results obtained with different algorithms.  
 287 We can see that for this particular problem, kofnGA, PSO, GA, and CE, managed  
 288 to obtain the  $-30.9263$  value for the objective function — the best solution known  
 289 so far (Wolters, 2015).



(a) The negative entropy as a function of algorithm, (lower is better).



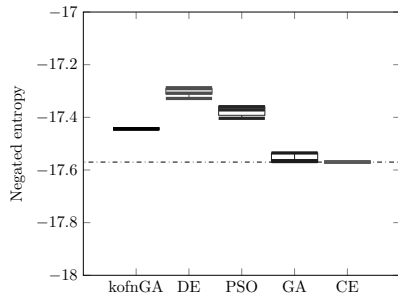
(b) CPU time as a function of algorithm.

Fig. 7: The  $m = 100$   $k = 15$  environmental monitoring design problem optimization results.

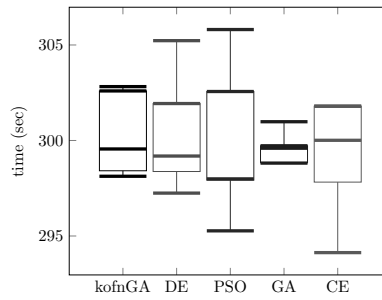
### 290 3.3 Custom covariance matrix

291 To test the performance of the proposed technique on a larger network design  
 292 problem, we consider a  $1000 \times 1000$  custom covariance matrix ( $m = 1000$ ), with  
 293  $k = 10$ ,  $k = 15$ , and  $k = 20$ . As mentioned in the beginning of this section, we

294 choose to examine a 1000 sites example, since many practical monitoring design  
 295 tasks generally contain several hundreds of potential locations for the placement of  
 296 monitoring stations. For this particular problem, we are not aware of any optimal  
 297 solution, since the full-enumeration procedure requires a consideration of a large  
 298 number of  $\binom{1000}{10} \approx 2.634 \times 10^{23}$ ,  $\binom{1000}{15} \approx 6.881 \times 10^{32}$ , and  $\binom{1000}{20} \approx 3.395 \times 10^{41}$   
 299 combinations, for  $k = 10$ ,  $k = 15$ , and  $k = 20$  cases, respectively. Therefore, we  
 300 simply compare all algorithms and deliver the best solution found. Fig. 8, Fig. 9,  
 301 and Fig. 10, summarize the obtained results. Among all algorithms, CE performed  
 302 in the best fashion. Namely, the CE method manages to find the smallest values of  
 303  $-17.5699$ ,  $-26.3084$ , and  $-35.0301$  of the objective function, for  $k = 10$ ,  $k = 15$ ,  
 304 and  $k = 20$ , respectively.

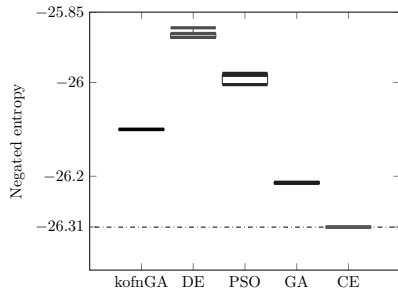


(a) The negative entropy as a function of algorithm, (lower is better).

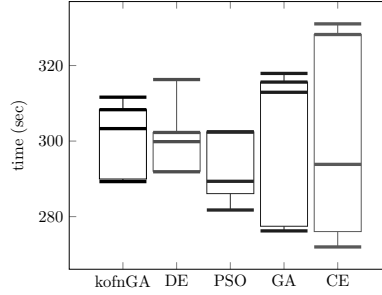


(b) CPU time as a function of algorithm.

Fig. 8: The  $m = 1000$   $k = 10$  custom monitoring design problem optimization results.

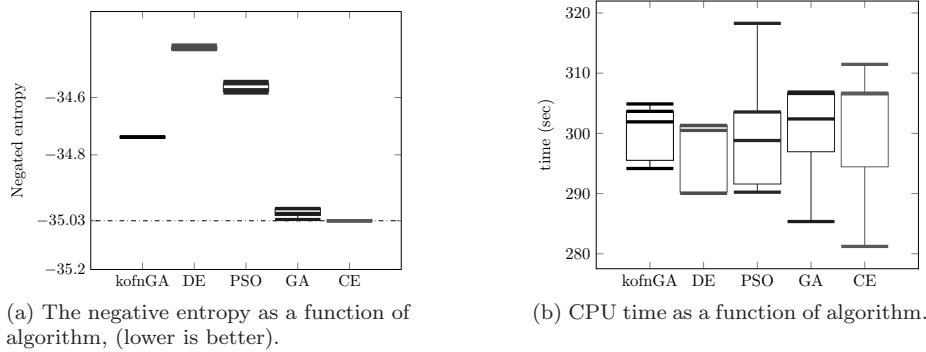


(a) The negative entropy as a function of algorithm, (lower is better).



(b) CPU time as a function of algorithm.

Fig. 9: The  $m = 1000$   $k = 15$  custom monitoring design problem optimization results.



(a) The negative entropy as a function of algorithm, (lower is better).

(b) CPU time as a function of algorithm.

Fig. 10: The  $m = 1000$   $k = 20$  custom monitoring design problem optimization results.

305 *Remark 1* Since only the CE algorithm obtained the best known objective function  
 306 values for all cases, we allowed all algorithms to run for an additional time. Specifi-  
 307 cally, we increased the CPU time limit by a factor of 10 for each global optimization  
 308 method. Each algorithm was executed for 5 times. As a consequence, both CE and  
 309 GA, managed to achieve the (possibly) optimal results for all  $k \in \{10, 15, 20\}$ . Ta-  
 310 ble 1 summarizes the best results obtained by all optimization methods when using  
 311 the the additional computational time.

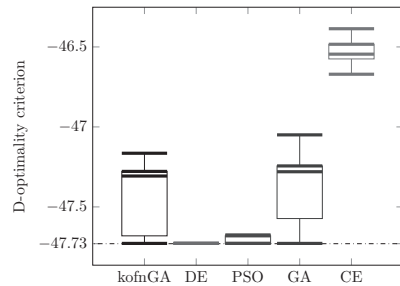
Table 1: The minimal negative entropy achieved by global optimization algorithms using the increased computational time limit.

algorithm	$k = 10$	$k = 15$	$k = 20$
kofnGA	-17.4434	-26.1002	-34.7382
DE	-17.4304	-26.1307	-34.6078
PSO	-17.3985	-26.0879	-34.5897
GA	<b>-17.5699</b>	<b>-26.3084</b>	<b>-35.0301</b>
CE	<b>-17.5699</b>	<b>-26.3084</b>	<b>-35.0301</b>

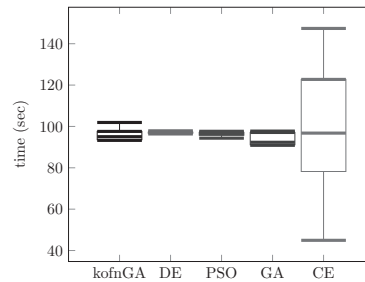
312 By analyzing the results from Sections 3.1, 3.2, and 3.3, one can arrive to  
 313 a wrong conclusion about the CE algorithm superiority as compared to other  
 314 methods. To see that there is no single best algorithm, we will examine a problem  
 315 from a different domain, the *D-optimal experimental design*.

## 316 3.4 D-optimal experimental design

317 In this section, we consider an instance of the D-optimal design problem called  
 318 the *robustness experiment* (Goos and Bradley, 2011). Following (Wolters, 2015),  
 319 we view this task as an OFSS problem with  $m = 81$  and  $k = 24$ . Consequently, a  
 320 full-enumeration procedure requires the evaluation of  $\binom{81}{24} \approx 2.305 \times 10^{20}$  combina-  
 321 tions. For this problem, Wolters (Wolters, 2015) manages to obtain the (probably)  
 322 optimal result (of  $-47.728$  on the negative log scale) in four out of 20 runs of the  
 323 problem using the kofnGA method. We repeat the experiment using DE, PSO,  
 324 GA, and CE algorithms. Similar to previous experiments, we set the parameters  
 325 such that the running times of all algorithms are comparable. Fig. 11 summarizes  
 326 the obtained results.



(a) The D-optimality as a function of algorithm, (lower is better).



(b) CPU time as a function algorithm.

Fig. 11: The  $m = 81$   $k = 24$  D-optimal design problem optimization results.

Table 2: Frequency of obtaining the best-known solution.

algorithm	success rate
kofnGA	15%
DE	100%
PSO	70%
GA	5%
CE	0%

327 Table 2 shows the frequencies of obtaining the best-known solution using dif-  
 328 ferent algorithms. It is interesting to note that the CE method always failed to  
 329 achieve the best known objective function value. Nevertheless, the DE algorithm  
 330 always managed to obtain the optimal result, and the PSO method performed rea-  
 331 sonably well, too. The kofnGA and GA methods achieved a relatively low success  
 332 rate of 15% and 5%, respectively.

## 4 Conclusion

In this study we showed that *any* global continuous optimization method is applicable for solving the optimal fixed-size subset selection problem. We presented a suitable embedding of the original discrete state space into the continuous state space, and introduced a transformation that allows to employ continuous optimization procedures for the original discrete problem solution. Our numerical study indicates that the achieved performance compares favorably with the existing open-source specifically designed software. We showed that by utilizing the proposed method, one can run several algorithms in parallel and report the best solution obtained. Moreover, this procedure can be carried out with a minimal development effort. Unsurprisingly, our numerical study indicates that there is no single best method available, and in practice, we advise to run several global optimization algorithms side-by-side.

While this study introduce a number of obvious benefits, it also suggests several promising directions for future research. First, we explored only a very limited number of continuous global optimization methods, namely, the Differential Evolution, the Particle Swarm Optimization, the Genetic Algorithm, and the Cross Entropy method. Our numerical study, and in particular, the D-optimal design example from the statistical domain, shows that it is important to explore additional continuous optimization techniques. The second promising direction is to explore different transformations from the continuous to the discrete space. A successful transformation should be constructed with a view to a satisfaction of smoothness condition of the objective function. A carefully designed embedding can assist a global optimization algorithm to improve both the speed of convergence and the solution obtained. Finally, it is important to note that the the Cross Entropy method deserves an additional attention as compared to other methods, since it can operate in stochastic environments. That is, it can be potentially used for an on-line monitoring network designs, and thus improve our capabilities of decision making under uncertainty.

## Declaration of interest statement

The author declares that there are no conflicts of interest in the preparation of this manuscript.

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