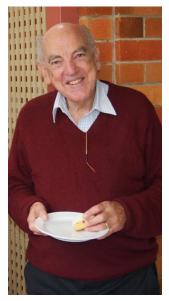
BRAVO for QED Queues

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Joint work with

Daryl J. Daley, The University of Melbourne, Johan van Leeuwaarden, EURANDOM, Eindhoven University of Technology.

Applied Probability Society Conference, Costa Rica, July, 2012.



Daryl J. Daley

$$Q(t) = Q(0) + \left(A(t) - L(t)\right) - \left(R(t) + D(t)\right)$$

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Daryl Daley, "*Queueing Output Processes*", Advances in Applied Probability, 1976.

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Some performance measures of interest

- The law of $\{D(t), t \ge 0\}$
- $\mathbb{E}[D(t)]$, Var(D(t))

•
$$\lambda^* := \lim_{t \to \infty} \frac{\mathbb{E}[D(t)]}{t}, \quad \overline{V} := \lim_{t \to \infty} \frac{\mathsf{Var}(D(t))}{t}, \quad \mathcal{D} := \frac{\overline{V}}{\lambda^*}$$

• Asymptotic normality: $D(t) \sim \mathcal{N}ig(\lambda^* t, \ \overline{V}tig)$, large t

- A (new) formula for asymptotic variance of outputs, $\mathcal{D} := \frac{V}{\lambda^*}$
- Single servers (older BRAVO results)
- Many server scaling (new BRAVO results)

Asymptotic Variance of Outputs

Finite Birth-Death Asymptotic Variance

- Irreducible birth-death process on finite state space
- Birth rates: $\lambda_0, \ldots, \lambda_{J-1}$
- Death rates: μ_1, \ldots, μ_J
- Stationary distribution: π_0, \ldots, π_J
- D(t) is number of downward transitions (deaths) during [0, t], each "filtered" independently with state-dependent probabilities, q₁,..., q_J.
- e.g. The departure process (served customers) in M/M/s/K+M systems

Of interest:

$$\mathcal{D} = rac{\overline{V}}{\lambda^*} = \lim_{t o \infty} rac{\mathsf{Var}ig(D(t)ig)}{\mathbb{E}[D(t)]}$$

Finite Birth-Death Asymptotic Variance Formula

Theorem: Daryl Daley, Johan van Leeuwaarden, Y.N. 2013

$$\mathcal{D} := \lim_{t \to \infty} \frac{\mathsf{Var}(D(t))}{\mathbb{E}[D(t)]} = 1 - 2\sum_{i=0}^{J} (P_i - \Lambda_i^*) \Big(q_{i+1} - \frac{\lambda^*}{\pi_i \lambda_i} (P_i - \Lambda_i^*) \Big),$$

with,

$$P_i := \sum_{j=0}^i \pi_j, \qquad \lambda^* := \sum_{j=1}^J \mu_j q_j \pi_j, \qquad \Lambda_i^* := \frac{\sum_{j=1}^i \mu_j q_j \pi_j}{\lambda^*}.$$

Note: In Weiss, Y.N. 2008, similar expression for case $q_i \equiv 1$

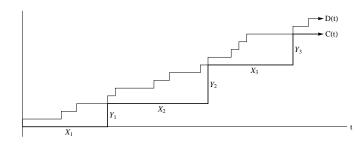
Note: In case $\lambda_i \equiv \lambda$, $q_i \equiv 1$:

$$\mathcal{D} = 1 - 2 \frac{\pi_J}{1 - \pi_J} \sum_{i=0}^{J} P_i \left(1 - \pi_J \frac{P_i}{\pi_i} \right)$$

Idea of Renewal Reward Derivation

"Embed" D(t) in a Renewal-Reward Process, C(t)

- (X_n, Y_n) \equiv (busy cycle, number served) in cycle *n*
- **2** $N(t) = \inf\{n : \sum_{i=1}^{n} X_i > t\}, \ C(t) = \sum_{i=1}^{N(t)} Y_i$
- Solution Asymptotic variance rates of C(t) and D(t) are equal
- 4 Known:
 - Asymptotic variance rate of C(t) is $\frac{1}{\mathbb{E}[X]} \operatorname{Var}(Y \frac{\mathbb{E}[Y]}{\mathbb{E}[X]}X)$
 - Systems of equations for 1'st, 2'nd and cross moments of X and Y



Single Server BRAVO (older results)

Here π_i is truncated geometric distribution when $\lambda \neq \mu$ and a uniform distribution when $\lambda = \mu$

Using
$$\mathcal{D} = 1 - 2 rac{\pi_J}{1 - \pi_J} \sum_{i=0}^J P_i \Big(1 - \pi_J rac{P_i}{\pi_i} \Big)$$
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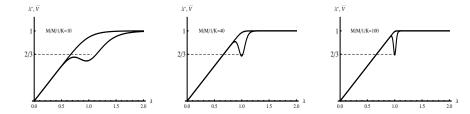
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$$\mathcal{D} = \left\{ egin{array}{cc} 1 + o_{\mathcal{K}}(1), & \lambda
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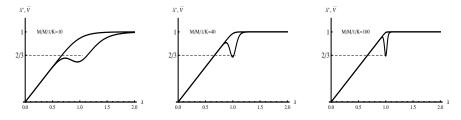
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We call this **BRAVO**:

Balancing Reduces Asymptotic Variance of Outputs

When $K = \infty$, the formula for \mathcal{D} does not hold. In this case,

$$\mathcal{D} = \begin{cases} 1, & \lambda \neq \mu, \\ ?, & \lambda = \mu. \end{cases}$$

A guess is $rac{2}{3}$, since for $K<\infty$, $\mathcal{D}=rac{2}{3}+o_{\mathcal{K}}(1)...$

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Theorem: Ahmad Al-Hanbali, Michel Mandjes, Y. N., Ward Whitt, 2011 For the M/M/1 queue with $\lambda = \mu$ and arbitrary initial conditions

of Q(0) (with finite second moments),

$$\mathcal{D} = 2\left(1 - \frac{2}{\pi}\right) \approx 0.727.$$

Proof based on analysis of classic Laplace transform of generating function of $D(\chi)$ where χ is an exponential random variable.

G/G/1 Queue

Moving away from the memory-less assumptions,

$$\mathcal{D} = \begin{cases} c_{\mathsf{a}}^2, & \lambda < \mu, \\ ?, & \lambda = \mu, \\ c_{\mathsf{s}}^2, & \lambda > \mu. \end{cases}$$

For M/M/1 it was $2(1-\frac{2}{\pi})...$

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Theorem:

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For the G/G/1 queue with $\lambda = \mu$, arbitrary finite second moment initial conditions (Q(0), V(0), U(0)), and finite fourth moments of the inter-arrival and service times,

$$\mathcal{D}=(c_a^2+c_s^2)\Big(1-\frac{2}{\pi}\Big).$$

Proof based on diffusion limit of $(D(n \cdot) - \lambda n \cdot)/\sqrt{\lambda n}$ as $n \to \infty$ (Iglehart and Whitt 1971). Fourth moments are a technical condition used in establishing uniform integrability.

G/G/1/K Queue

$$\mathcal{D} = \left\{ egin{array}{ll} c_{s}^{2}+o_{\mathcal{K}}(1), & \lambda < \mu, \ ?, & \lambda = \mu, \ c_{s}^{2}+o_{\mathcal{K}}(1), & \lambda < \mu. \end{array}
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For M/M/1/K it was $\frac{2}{3} + o_K(1)$, for G/G/1 it was $(c_a^2 + c_s^2)(1 - \frac{2}{\pi})...$

G/G/1/K Queue

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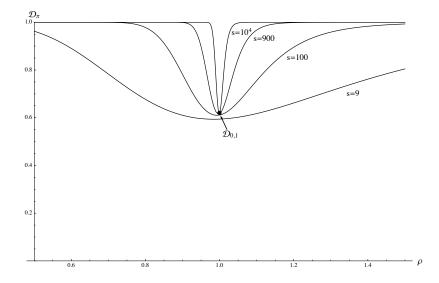
Conjecture (numerically tested), Y.N., 2011

For the G/G/1/K queue with $\lambda = \mu$ and arbitrary initial conditions and light-tailed service and inter-arrival times,

$$\mathcal{D}=(c_a^2+c_s^2)\frac{1}{3}+O(\frac{1}{K}).$$

Numerical verification done by representing the system as PH/PH/1/K MAPs

Many Servers



Quality and Efficiency Driven (QED) Scaling Regime

A sequence of systems

Consider a sequence of M/M/s/K queues with increasing s = 1, 2, ... and with $\rho_s := \frac{\lambda}{s\mu}$ and K_s such that,

$$(1 - \rho_s)\sqrt{s} \to \beta \in (-\infty, \infty)$$

 $\frac{K_s}{\sqrt{s}} \to \eta \in (0, \infty)$

So for large s:

$$ho_{s}pprox 1-eta/\sqrt{s}$$

 $K_{s}pprox \eta\sqrt{s}$

Halfin, Whitt, 1981, Garnett, Mandelbaum, Reiman 2002, Borst, Mandelbaum, Reiman, 2004, Whitt, 2004, Pang, Talreja, Whitt, 2007, Janssen, van Leeuwaarden, Zwart, 2011, Kaspi, Ramanan 2011, first session of this morning....

- $\bullet\,$ Probability of delay converges to a value \in (0,1)
- Mean waiting times are typically $O(s^{-1/2})$
- Large queue lengths almost never occur
- Quick mixing times
- In applications: Call-centers (etc...) describes behavior well and allows for asymptotic approximate optimization of staffing etc...
- How about BRAVO?

BRAVO for QED Queues

Theorem: Daryl Daley, Johan van Leeuwaarden, Y.N. 2013

Consider QED scaling with $\beta \neq 0$:

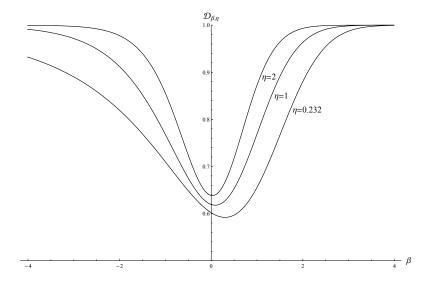
$$\mathcal{D}_{eta,\eta}:=\lim_{s,K o\infty}\lim_{t o\infty}rac{Varig(D(t)ig)}{\mathbb{E}ig(D(t)ig)},$$

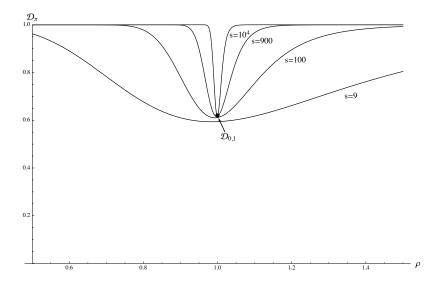
$$\mathcal{D}_{\beta,\eta} = 1 - \frac{2\beta^2 e^{-\beta\eta} h^2}{\phi(\beta)} \int_{-\beta}^{\infty} \left(1 - \beta e^{-\beta\eta} h \frac{\Phi(-u)}{\phi(u)} \right) \Phi(-u) \, du$$
$$+ 2e^{-\beta\eta} h (1 + e^{-\beta\eta} h) \left(1 - \beta\eta - e^{-\beta\eta} + (1 - 2\beta\eta e^{-\beta\eta} - e^{-2\beta\eta}) h \right)$$

where

$$h = \lim_{s \to \infty} \frac{\mathbb{P}(Q_s \ge s)}{1 - e^{-\beta\eta}} = \frac{1}{1 - e^{-\beta\eta} + \frac{\beta\Phi(\beta)}{\phi(\beta)}}$$

BRAVO viewed through the QED lens





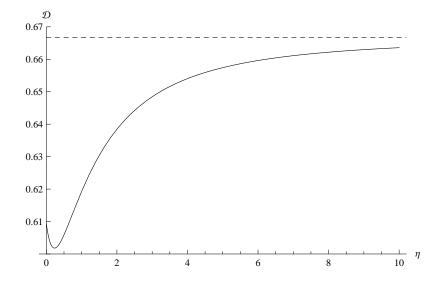
M/M/s/K QED BRAVO with $\rho \equiv 1 \ (\beta = 0)$

Theorem: Daryl Daley, Johan van Leeuwaarden, Y.N. 2013

$$\frac{K_s}{\sqrt{s}} \to \eta \in (0,\infty)$$

$$\mathcal{D}_{0,\eta} \coloneqq \lim_{s, K \to \infty} \lim_{t \to \infty} rac{Var(D(t))}{\mathbb{E}(D(t))},$$
 $\mathcal{D}_{0,\eta} = rac{2}{3} - rac{\left(6 - rac{3\pi}{2}
ight)\eta - rac{1}{2}\pi\sqrt{rac{\pi}{2}} + 3\sqrt{2\pi}(1 - \log 2)}{3\left(\eta + \sqrt{rac{\pi}{2}}
ight)^3}.$

$\mathsf{M}/\mathsf{M}/s/\lfloor\eta\sqrt{s} floor$ $s o\infty$ at $ho\equiv 1~(eta=0)$



Summary

Known BRAVO constants:

- Single server finite buffer: 2/3
 - (for G/G replace 2 by $c_a^2 + c_s^2$)
- Single server infinite buffer $2(1-2/\pi)$: (for G/G replace 2 by $c_a^2 + c_s^2$)
- Memoryless many servers finite buffer: $\mathcal{D}_{0,\eta} \in [0.6, 2/3]$

Not yet known:

- Memoryless many servers infinite buffer.
- Many servers without memoryless assumptions
- Systems with reneging or other customer loss mechanisms

Other questions: How can BRAVO be harnessed in practice? Why does BRAVO occur?

- Daryl J. Daley, Johan van Leeuwaarden and Y.N., "BRAVO for QED Finite Birth-Death Queues", preprint.
- Daryl Daley, "*Revisiting queueing output processes: a point process viewpoint*", Queueing Systems, 68, pp. 395–405, 2011.
- Y.N., "The variance of departure processes: puzzling behavior and open problems", Queueing Systems, 68, pp. 385–394, 2011.
- Ahmad Al-Hanbali, Michel Mandjes, Y.N. and Ward Whitt, "*The asymptotic variance of departures in critically loaded queues*", Advances in Applied Probability, 43, pp. 243–263, 2011.
- Y.N. and Gideon Weiss, "*The asymptotic variance rate of the output process of finite capacity birth-death queues*", Queueing Systems, 59, pp. 135–156, 2008.